

C455 Algorithms Analysis Recurrence Relations

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Definitions

- A **recurrence relation** for a sequence $F(n)$, $n \in \mathbb{N}$, is an equation of the form
- $F(n) = \Phi(n, F(n-1), F(n-2), \dots, F(0))$
- A solution for the recurrence relation is a formula for $F(n)$ that only depends on n .
- A **linear** recurrence relation is of the form
- $F(n) = c_1(n)F(n-1) + c_2(n)F(n-2) + \dots + c_n(n)F(0) + \phi(n)$
- If $\phi(n)$ is missing, then the relation is called **homogeneous**.

Definitions

- A **first order** recurrence relation depends only on n and $F(n-1)$.
- A **second order** relation depends on n , $F(n-1)$, $F(n-2)$. Third order is defined similarly and so on...
- For a recurrence relation of first order we must know $F(0)$ (or $F(1)$ or wherever the sequence starts from).
- For a recurrence relation of second order we need to know $F(0)$ and $F(1)$ and so on.

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Computational Efficiency

```
■  $F(1)=1, F(2)=2, F(n)=F(n-2)+nF(n-1)$   
long F(long n)  
{  
    if (n<=2)  
        return n;  
    else  
        return F(n-2)+n*F(n-1);  
}  
// Exponential complexity.  
// Lots of repetition
```

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```

long F(long n)
{
    if (n <= 2)
        return n;
    long F_n=2, F_n_1=1,
        F_n_2;
    for (i=3; i<=n; ++i) {
        F_n_2 = F_n_1;
        F_n_1 = F_n;
        F_n = F_n_2 - n*F_n_1;
    }
    return F_n;
}

```

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Examples

- **A5 a)** $F(n) = \frac{n^2 + 1g(n)}{\sqrt{n+1}}$ F(1), F(8), F(t), F(n-1), F(\sqrt{n}).
- **b)** $F(n) = (F(n-1))^2 - 3 F(n-3)$
- F(t), F(n-1), F(n+3), F($\lfloor n/2 \rfloor$)
- **A7)** $F(0)=2$, $F(n) = (F(n-1))^2 - 3 F(n-1)$
- Find F(4).
- **A9 a)** $F(n) = F(n-1) + 1/(n(n+1))$
- Verify the solution: $F(n) = \alpha + n/(n+1)$
- **b)** $F(n) = 4F(n-1) + 12F(n-2)$
- Solution: $F(n) = \alpha(-2)^n + \beta 6^n$

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Examples

- **A11)** Find solutions for
- **a)** $F(0)=1, F(n)=n F(n-1)$
- **b)** $F(0)=0, F(1)=-1,$
 $F(n) = -2F(n-1)-F(n-2)$
- **A13)** Order, linear, homogeneous?
- **a)** $F(n)=(F(n-1))^2-3F(n-3)$
- **b)** $F(n)=F(n-1)+2F(n-2)+3 (2)^{n-1}$
- **c)** $F(n)=(F(n-1)+F(n-2))/2$
- **d)** $F(n)=(F(n-1)+1)/2$

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More Examples

- **A13f)**
$$\frac{2^n F(n) - 2^{n+1} F(n-1)}{F(n-2) + F(n-3)} = 1 + \lfloor \lg(n!+1) \rfloor$$
- **A15.** Given the function F from ex. A3, what's wrong with the following code:

```
for (int k=1; k<=10000; ++k)
    cout << "F (" << k <<")="
        << F(k) << endl;
```

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Linear Homogeneous RR with Constant Coefficients

- **Theorem B1 (Superposition).** Suppose that $F(n)=g(n)$ and $F(n)=h(n)$ are two solutions for
- $F(n)=c_1(n)F(n-1)+c_2(n)F(n-2)+ \dots + c_n(n)F(0)$
- Then every linear combination of g and h is also a solution to the RR:
- $F(n) = \alpha g(n) + \beta h(n)$
- Also true for $\alpha = \beta = 0$

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General Solutions

- **Theorem B2.** The solution of the recurrence relation
- $F(n)=c_1 F(n-1), F(n_0)=a$
- is the following:
- $F(n)=\alpha c_1^n$ for all $n \geq n_0$
- **Theorem B3.** The solution for the RR
- $F(n)=c_1 F(n-1)+c_2 F(n-2)$
- is of the form
- $a_1(r_1)^n+a_2(r_2)^n$ or $(a_1+a_2n)(r_1)^n$
- where r_1 and r_2 are the solutions of
- $r^2=c_1r+c_2$

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Examples

- **B4.** Find the solution for:
 - $F(n) = -F(n-1) + 2F(n-2)$, $F(0) = 3$, $F(1) = 5$
- **B5.** Give the form of the solution for
 - a) $F(n) = 5F(n-1)$
 - b) $F(n) = 2F(n-1)/3$
- **B7.** Solve the following:
 - a) $F(0) = 1$, $F(n) = 2F(n-1)$
 - d) $T(0) = 3$, $T(n) = T(n-1)/2$

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Examples 2nd Order

- **B5.** Give the form of the solution for
 - d) $F(n) = F(n-1) + 2F(n-2)$
 - f) $F(n) = -3F(n-1) + F(n-2)$
- **B9.** Solve the following
 - a) $F(1) = 2$, $F(2) = 0$, $F(n) = F(n-1) + 6F(n-2)$
 - c) $f(0) = 0$, $f(1) = 1$, $f(n) = f(n-1) + f(n-2)$
(Fibonacci numbers)
 - d) $A(0) = -3$, $A(1) = 3$, $A(j) = 6A(j-1) - 9A(j-2)$

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Third Order RR

- The solution for a RR of the form
- $F(n) = c_1 F(n-1) + c_2 F(n-2) + c_3 F(n-3)$
- is of one of the forms
- $F(n) = a_1(r_1)^n + a_2(r_2)^n + a_3(r_3)^n$
- $F(n) = (a_1 + a_2n)(r_1)^n + a_3(r_3)^n$
- $F(n) = (a_1 + a_2n + a_3n^2)(r_3)^n$
- where r_1 , r_2 , and r_3 are the solutions for
- $r^3 = c_1r^2 + c_2r + c_3$

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Examples 3rd Order

- $F(n) = 3 F(n-1) - 4 F(n-3)$
- $F(0) = 4$
- $F(1) = 4$
- $F(2) = 34$

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Linear Nonhomogeneous RR

- **Theorem C1.** Suppose $F(n)=g(n)$ is a solution for the relation
- $F(n)=c_1(n)F(n-1)+c_2(n)F(n-2) + \dots c_n(n)F(0) + \phi(n)$
- and suppose that $h(n)$ is a solution for
- $F(n)=c_1(n)F(n-1)+c_2(n)F(n-2) + \dots c_n(n)F(0)$
- Then $F(n) = g(n) + \alpha h(n)$ if also a solution for the first RR.

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Theorem C2

- Suppose $F(n)=g(n)$ and $F(n)=h(n)$ are solutions for the relations
- $F(n)=c_1(n)F(n-1)+c_2(n)F(n-2) + \dots c_n(n)F(0) + \phi(n)$
- $F(n)=c_1(n)F(n-1)+c_2(n)F(n-2) + \dots c_n(n)F(0) + \psi(n)$
- Then $F(n)=\alpha g(n) + \beta h(n)$ is a solution for
- $F(n)=c_1(n)F(n-1)+c_2(n)F(n-2) + \dots c_n(n)F(0) + \alpha \phi(n) + \beta \psi(n)$

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Theorem C3

- Suppose $F(n)=g(n)$ and $F(n)=h(n)$ are solutions for the same RR
- $F(n)=c_1(n)F(n-1)+c_2(n)F(n-2) + \dots + c_n(n)F(0) + \phi(n)$
- The $F(n)=g(n)-h(n)$ is a solution for the homogeneous RR
- $F(n)=c_1(n)F(n-1)+c_2(n)F(n-2) + \dots + c_n(n)F(0)$

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Example

- $F(n)=3F(n-1)+4(5)^n+7$
- Decompose as
- $F(n)=3F(n-1)+4(5)^n$
- $F(n)=3F(n-1)+7$
- Solutions of the form
- $F(n)=\beta (5)^n$ and $F(n)=\gamma$ (constant)
- Use Theorems C2 and C3.

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General Procedure

- a) Find the most general solution for the homogeneous part of the RR
- b) Find a particular solution for the nonhomogeneous RR
- c) Add the two to obtain a general solution for the RR
- d) Use the initial values for the sequence to determine the values of all the constants

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Polynomial NonH RR First Order

- $F(n) = c_1 F(n-1) + P(n)$, where $P(n)$ is a polynomial.
- Solution:
 - if $c_1 \neq 1$, $F(n) = Q(n)$, where $Q(n)$ is a polynomial of the same degree as $P(n)$
 - if $c_1 = 1$, $F(n) = Q(n)$, where $Q(n)$ is a polynomial of a degree 1 greater than P with the constant term being 0.
- Examples: a) $F(n) = 2F(n-1) + 2n^3 - 5n$
- b) $F(n) = F(n-1) + 3n$ c) $F(n) = F(n-1) + 7$

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Second Order

- $F(n)=c_1F(n-1)+c_2F(n-2)+P(n)$, where P is a polynomial.
- Solution: let r_1 and r_2 be the solutions for $r^2=c_1r+c_2$.
- If $r_1 \neq r_2$, then the solution is $Q(n)$, a polynomial of the same degree as P .
- If one of r_1 and r_2 is 1, then the solution is $Q(n)$, a polynomial of degree 1 greater than P , constant term in Q being 0.
- If $r_1=r_2=1$, then the solution is $Q(n)$, a polynomial of degree 2 greater than P , constant term =0 and coefficient of $n=0$.

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Examples

- **C8**
- a) $F(n)=4F(n-1)+5F(n-2)+2n$
- b) $F(n)=-5F(n-1)+6F(n-2)+2n$
- c) $F(n)=2F(n-1)-F(n-2)+2n$
- **C9.** $F(0)=1, F(n)=-2F(n-1)+n-1$
- **C11.** To compute $F(n)=\sum k^2$ for k from 1 to n , we can express it as
- $F(n) = F(n-1) + n^2$

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Theorem C10 (Matching Coefficients)

- Suppose we have an equation of the form
- $a_k n^k + a_{k-1} n^{k-1} + \dots + a_1 n + a_0 = b_k n^k + b_{k-1} n^{k-1} + \dots + b_1 n + b_0$
- and suppose that the equation is satisfied by an infinite number of values of n . Then $a_0 = b_0, a_1 = b_1, \dots, a_k = b_k$.
- Similarly, the equation
- $c_1 q_1^n + c_2 q_2^n + \dots + c_k q_k^n = d_1 q_1^n + d_2 q_2^n + \dots + d_k q_k^n$
- is satisfied by an infinite number of values of n iff all coefficient c_i and d_i are respectively equal.

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Exponential NonH Linear RR with Constant Coefficients

- **First Order.** $F(n) = c_1 F(n-1) + d q^n$
- If $c_1 \neq q$, $F(n) = \gamma q^n$
- If $c_1 = q$, $F(n) = \gamma n q^n$
- **Examples.**
- C13 a) $F(n) = F(n-1) + 3 \cdot 2^n$
- b) $F(n) = 2F(n-1) + 2 \cdot 3^n + 3 \cdot 2^n$

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Second Order

- $F(n)=c_1F(n-1)+c_2F(n-2)+dq^n$
- Let r_1 and r_2 be the solutions for $r^2=c_1r+c_2$.
- If $r_1 \neq r_2$, then $F(n)=\gamma q^n$
- If one of r_1 and r_2 is equal to q , then $F(n)=\gamma nq^n$
- If both $r_1=r_2=q$, then $F(n)=\gamma n^2q^n$
- *Note1:* only one value of γ will satisfy the recurrence relation. It is determined by replacing the solution in the RR.
- *Note2:* We still need to add to this the general solution to the homogeneous part of the equation. The coefficients of that part are determined by $F(0)$, $F(1)$, etc.

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Examples

- **C14 a)** $F(n)=5F(n-1)-6F(n-2) + (-1)^n$
- **b)** $F(n)=F(n-1)+2F(n-2) + (-1)^n$
- **c)** $F(n)=6F(n-1)-9F(n-2)+1/2^n+3^n$
- **C15.** $F(1)=1$, $F(2)=2$,
- $F(n)=2F(n-1)+4F(n-2)-5n+2(-1)^{n+1}$

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C17

- Which of these are polynomials and of what degree:
- a) n^3+3n^2-2n+1 of n
- b) 3^n-2^n+1 of n
- c) 2^{10} of n
- d) $\sqrt{4n^2-3n+7}$ of n
- e) $(k+1)/2$ of k
- f) $\lfloor (k+1)/2 \rfloor$ of k
- g) $p \lg m$ of m (p is constant)
- h) $p \lg m$ of p (m is constant)

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C19

- a) $F(n)=3F(n-1)+3n^2-2n+1$
- b) $F(n)=F(n-1)+3n^2-2n+1$
- c) $G(k)=G(k-1)-2$
- d) $T(n)=(T(n-1)+n)/3$
- e) $F(n)=3F(n-1)-2F(n-2)-(n+1)/5$
- f) $g(n)=-2g(n-1)+5g(n-2)-4n^2+6$
- g) $F(n)=2F(n-1)-F(n-2)-n^3$
- h) $M(k)=5M(k-1)-6M(k-2)+k^2$

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Recurrence Relations in $\lfloor n/2 \rfloor$ and $\lceil n/2 \rceil$

- **D1.** $F(1)=a, F(n) = b F(\lfloor n/2 \rfloor)$
- $F(n)=b^2 F(\lfloor n/4 \rfloor)=\dots b^k F(\lfloor n/2^k \rfloor)$
- Iterate until $\lfloor n/2^k \rfloor = 1, k = \lfloor \lg n \rfloor$
- $F(n) = a b^{\lfloor \lg n \rfloor}$
- **D2.** $F(1)=a, F(n)= b F(\lfloor n/2 \rfloor) + 1$
- Same procedure
- $F(n) = a b^{\lfloor \lg n \rfloor} + (b^{\lfloor \lg n \rfloor} - 1)/(b - 1)$
- $b=1 \Rightarrow F(n)=a + \lfloor \lg n \rfloor$

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More Examples

- **D3.** $F(1)=a, F(n)= b F(\lfloor n/2 \rfloor) + n$
- Same kind of iteration.
- $F(n)=a b^{\lfloor \lg n \rfloor} + \sum_{k=1}^{\lfloor \lg n \rfloor} \lfloor n/2^k \rfloor b^k$
- $F(n)=\Theta(n^{\lg b})$
- **D4.** $F(1)=a, F(n)=F(\lfloor n/2 \rfloor) + F(\lceil n/2 \rceil)$
- $F(n)=an$
- **D5.** $F(1)=a, F(n)=F(\lfloor n/2 \rfloor)+F(\lceil n/2 \rceil)+1$
- $F(n)=(a+1)n-1$

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Examples

- **D6.** $F(1)=a,$
 $F(n) = F(\lfloor n/2 \rfloor) + F(\lceil n/2 \rceil) + n$
- $F(n)=(a+2)n+n \lfloor \lg n \rfloor - 2^{\lfloor \lg n \rfloor + 1}$
- $F(n) \sim n \lg n$
- **D7.** $F(0)=1$
 $F(n)=2F(\lfloor n/2 \rfloor) - 7 \lfloor n/2 \rfloor + 5n+3$
- $F(n)=\Theta(n \log n)$

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Other RRs

1. Consult a reference book.
2. Calculate a good number of terms and try to "guess" the general formula that you can then verify.
3. Manipulate the equation to obtain a more easily solvable one.
4. Transform the equation into an inequality that can be more easily solved.

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Examples

- **E1.** $F(0)=a, F(n)=F(n-1)+\phi(n)$
- **E2.** $F(1)=1, nF(n)=(n-1)F(n-1)+2^n$
- **E3.** $F(1)=1, nF(n)=(n+1)F(n-1)+1$
- **E4.** $F(0)=5, F(n)=F(\lfloor (n-1)/2 \rfloor)+3$
- **E5.** $F(1)=2, F(n)=F(\lfloor n/2 \rfloor) + n \lfloor n/2 \rfloor$
- **E6.** $F(0)=1, F(n)=n/(n+2)F(n-1)+\lfloor n/2 \rfloor$

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Theorem E7

- Suppose $F(0), F(1), F(2), \dots$ satisfy a recurrence inequality of the form $F(n) < c(n)F(n-1) + \phi(n)$, where $c(n) > 0$ for all $n > 0$. Let $G(n)$ be defined by $G(0) = F(0), G(n) = c(n)F(n-1) + \phi(n)$ for $n > 0$. Then $F(n) < G(n)$ for any $n > 0$.

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More Examples

- **E9.** $F(n) = b F(\lceil (n+1)/2 \rceil) + \phi(n)$. Prove that it can be transformed into $F(n) = G(n-1)$ where $G(n) = b G(\lfloor n/2 \rfloor) + \phi(n+1)$
- **E11.** Solve the problem $F(2) = 5$, $F(n) = 2F(\lfloor n/2 \rfloor + 1) + 3 \lfloor n/2 \rfloor$, $n \geq 3$. Substitute $F(n) = G(n-1)$.
- **E15.** $F(1) = 5$, $F(n) = F(\lfloor n/2 \rfloor) + \sqrt{n}$